

# Verification of Multi-Agent Properties in Electronic Voting: A Case Study

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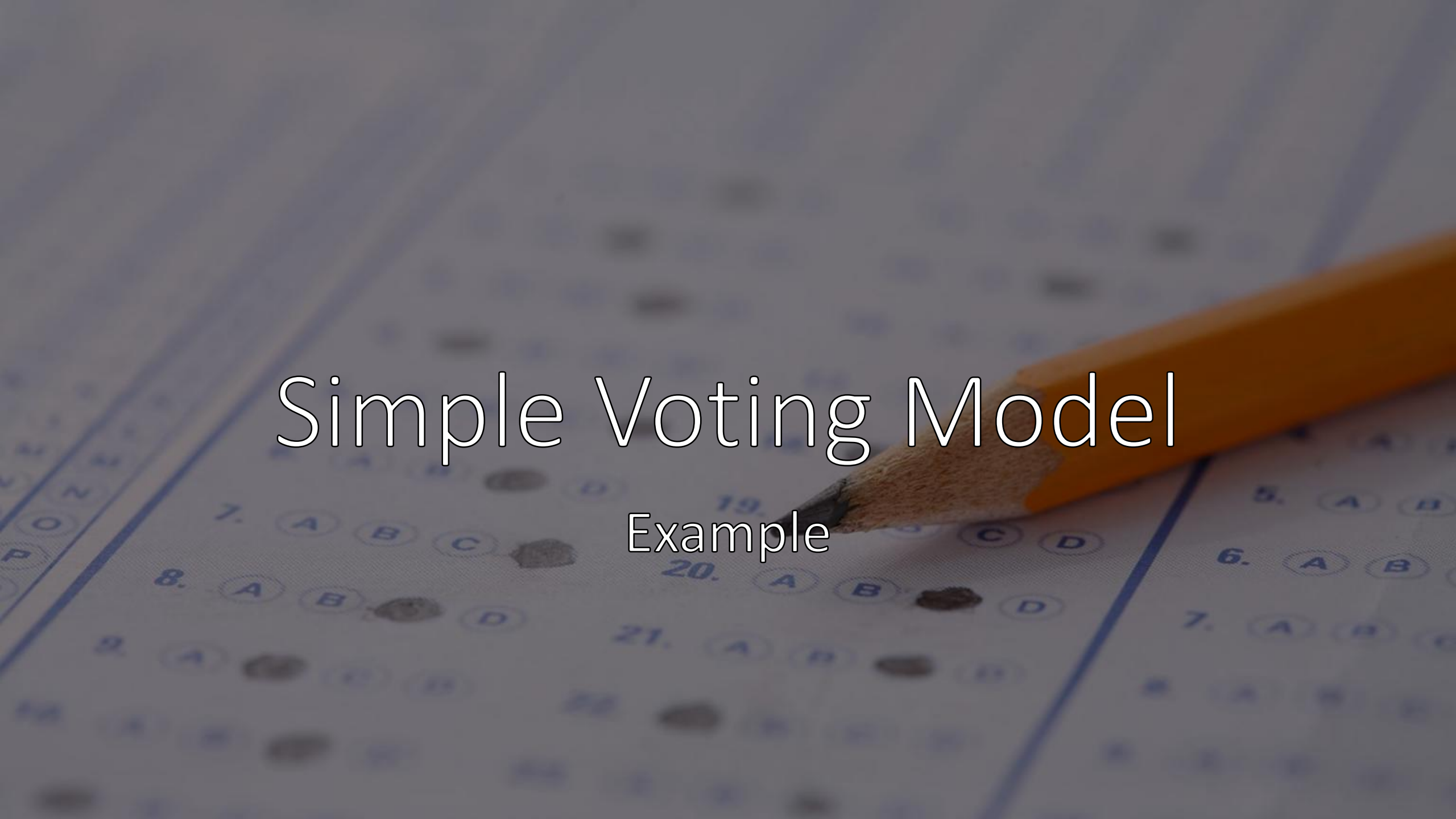
# The problem

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- Verification of strategic abilities under **imperfect information**
- Logic: **ATL<sub>ir</sub>**
- Complexity:  $\Delta_2^P$  – complete

# Simple Voting Model

Example



# Agents

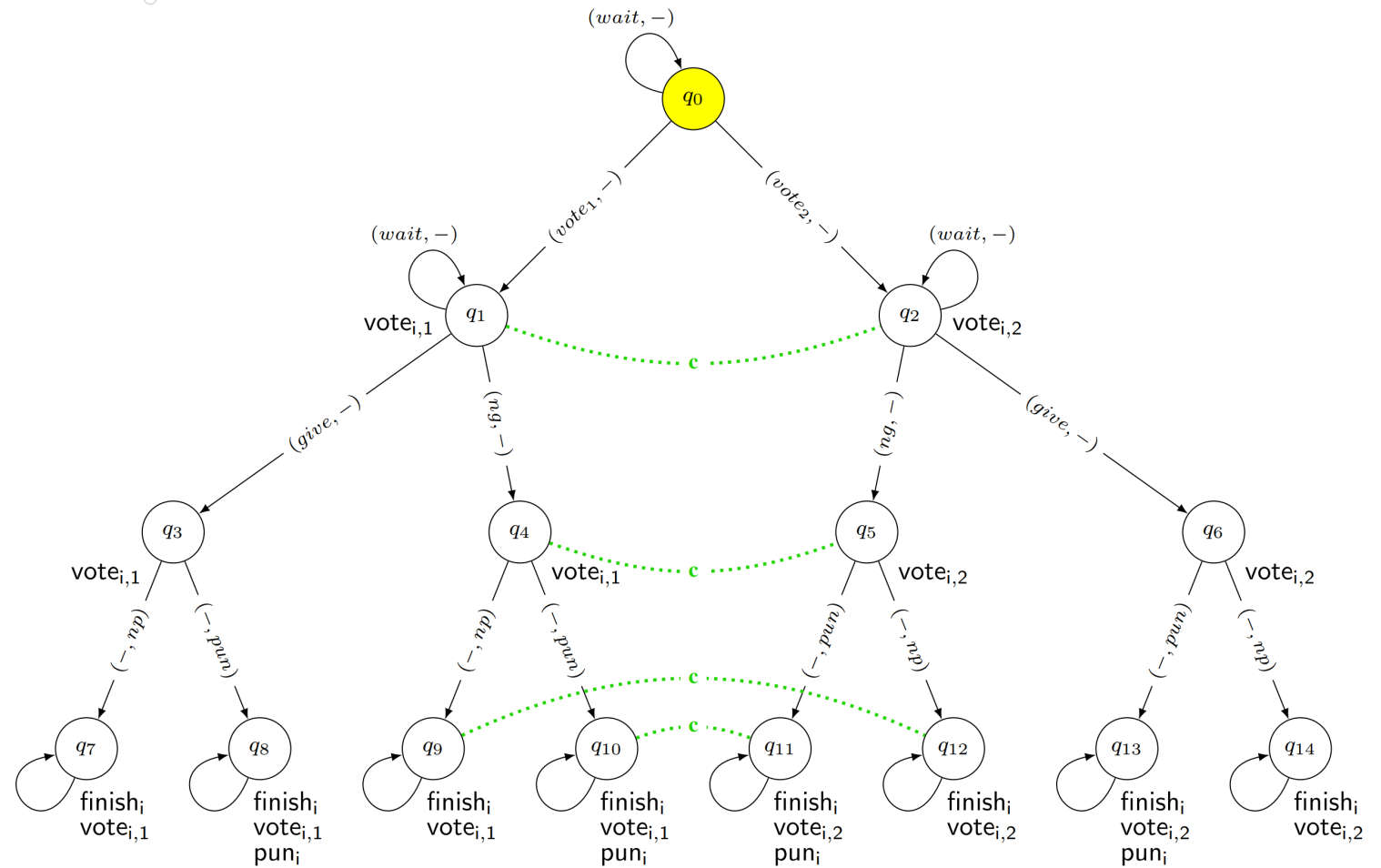
## Voter

- Casts her vote
- Decides to show (or not) her vote to the Coercer

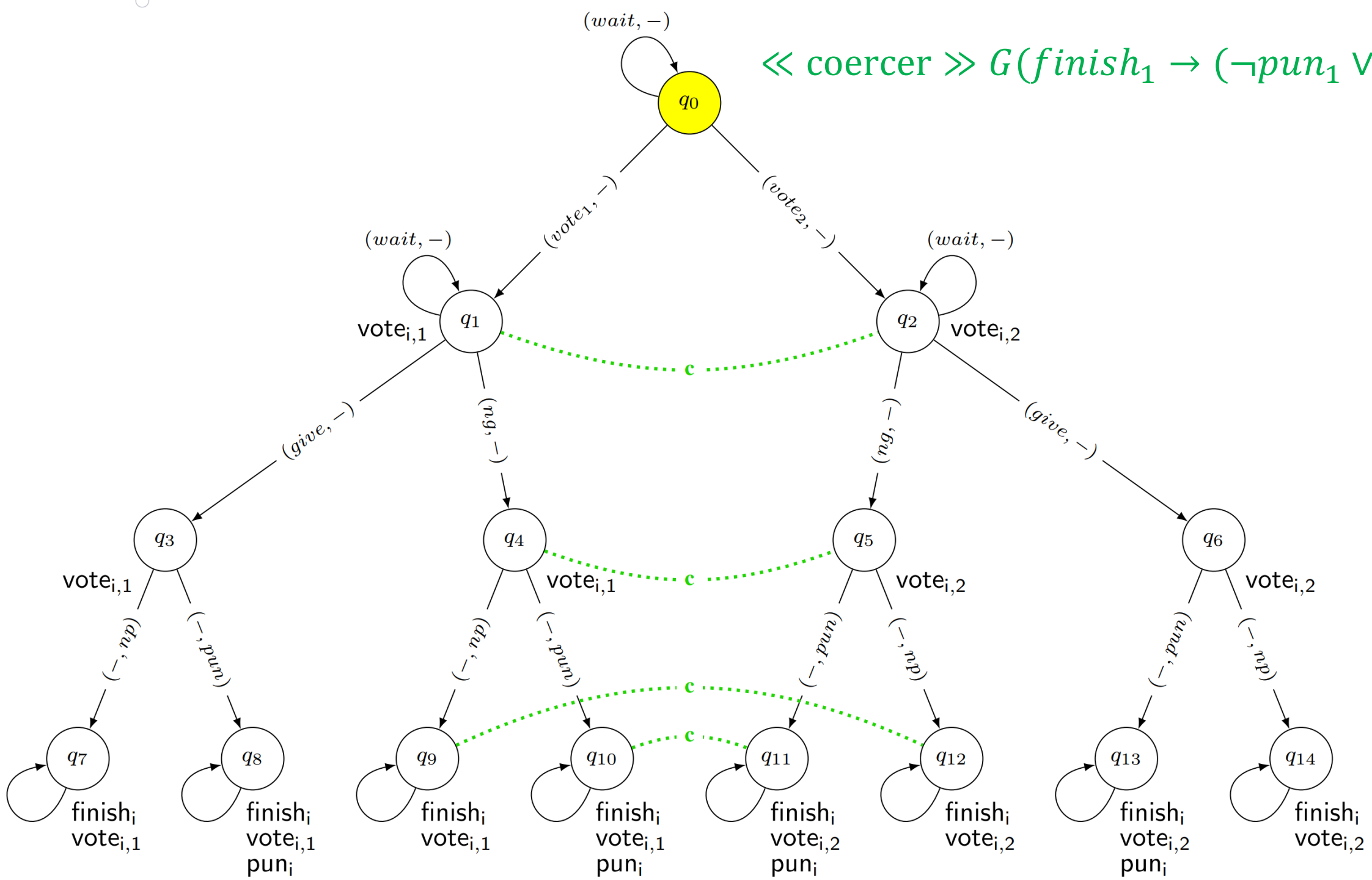
## Coercer

- Gets (or not) the vote from the Voter
- Decides to punish (or not) the Voter

1 Voter  
1 Coercer  
2 Candidates

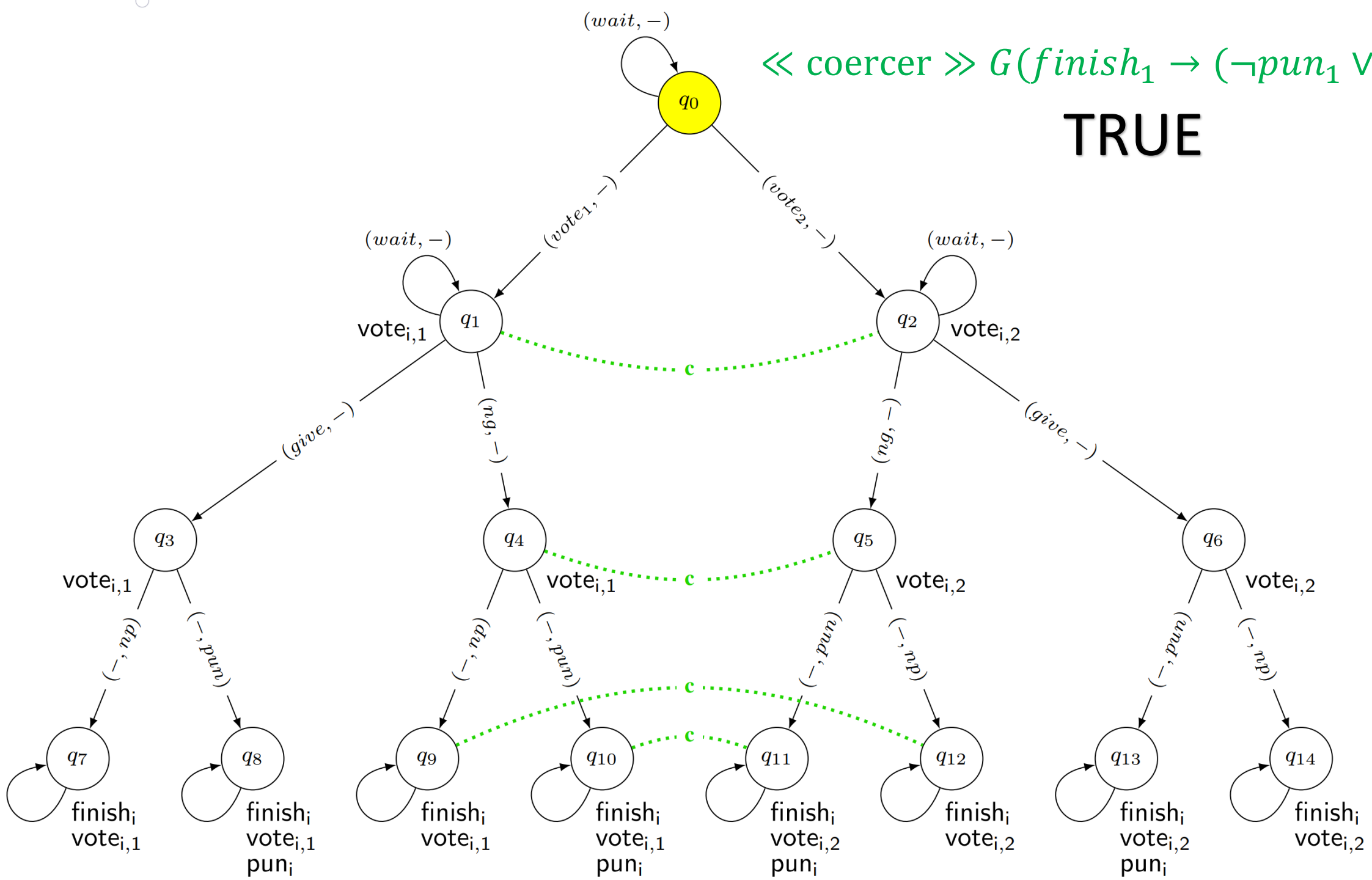


$\ll \text{coercer} \gg G(\text{finish}_1 \rightarrow (\neg \text{pun}_1 \vee \text{vote}_{1,1}))$

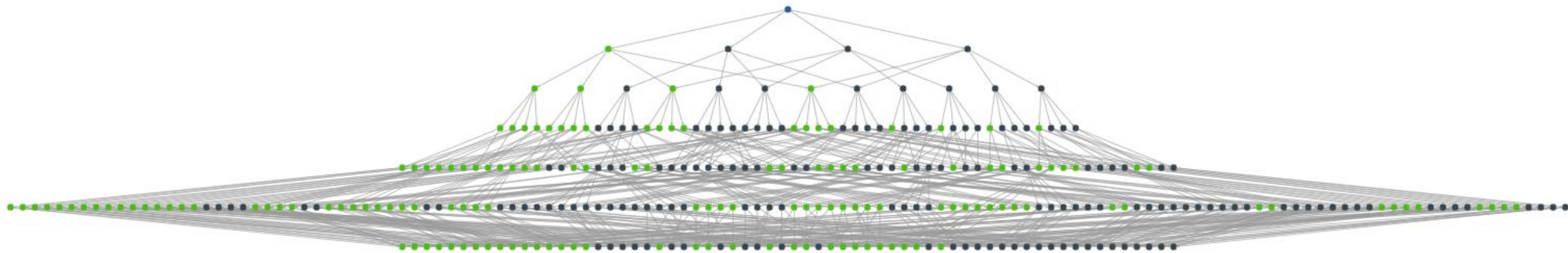


$\ll \text{coercer} \gg G(\text{finish}_1 \rightarrow (\neg \text{pun}_1 \vee \text{vote}_{1,1}))$

TRUE



# 2 Voters, 1 Coercer, 2 Candidates





# The solution(?)



Fixpoint approximations



DFS and DominoDFS strategy synthesis



Parallel DFS strategy synthesis



Partial-order reductions

## Fixpoint approximations

- Fixpoint computation is (usually) efficient
- Fixpoint equivalences do not hold for  $ATL_{ir}$
- **Lower bound:** translation to  $AE\mu C$
- **Upper bound:**  $ATL_{Ir}$  (perfect information)
- Sometimes bounds don't match

## DFS strategy synthesis

- Recursive search from the initial state
- Synthesize winning strategy during the search
- Better than exhaustive search through the entire strategy space
- Handling epistemic classes can be troublesome

# DominoDFS strategy synthesis

- DFS + domination relations
- Observation: some strategies **dominate** others
- Dominated strategies can be omitted during the search

# Parallel DFS strategy synthesis

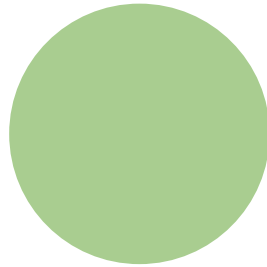
- Main problems to consider:
  - It is difficult (if not impossible) to split the model data between processes
  - Epistemic classes can join states in different parts of the model
  - Backtracing is not as simple as it seems
- Several different approaches to parallelization
- Best promising approach:
  - Split the work early (preferably from the initial state)
  - Each process has own copy of the whole model
  - Split by agent-controlled transitions

## Partial-order reductions

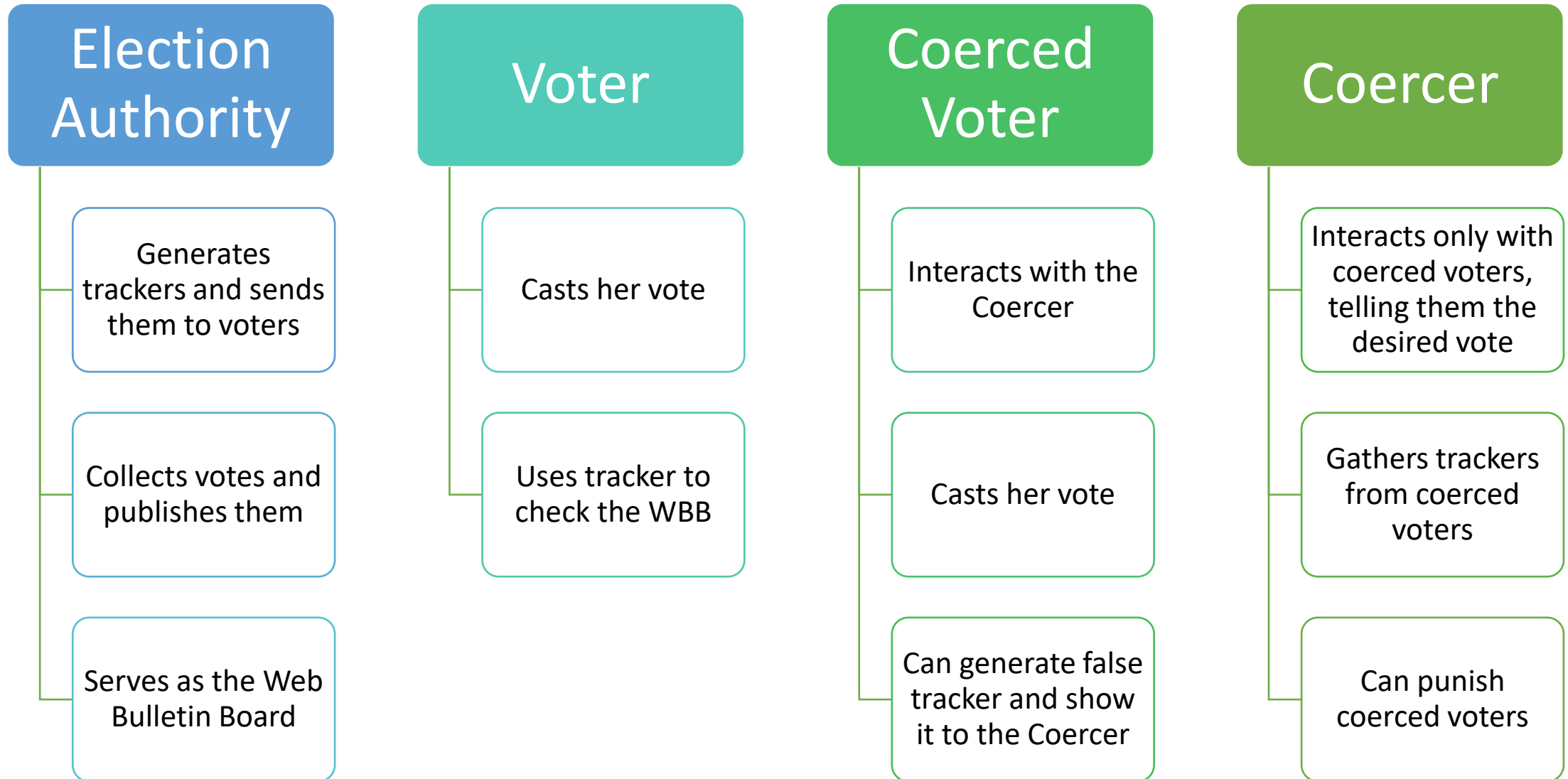
- Asynchronous models
- State-space explosion related to interlacing
- Effective reduction methods exists for LTL and can be adapted to  $ATL_{ir}$

# Selene e-voting Protocol Model

Case Study



# Agents





# Re-voting scheme

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Coerced voter can vote several times

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Each vote, apart from the last one, is shared with the coercer

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**Last vote** (if cast) **is private**

# Coerced Voter (3 candidates, 3 revotes)

```
Agent VoterC[1]:
init start
shared coerce1_aID: start -> coerced [aID_required=1]
shared coerce2_aID: start -> coerced [aID_required=2]
shared coerce3_aID: start -> coerced [aID_required=3]
select_vote1: coerced -> prepared [aID_vote=1, aID_prep_vote=1]
select_vote2: coerced -> prepared [aID_vote=2, aID_prep_vote=2]
select_vote3: coerced -> prepared [aID_vote=3, aID_prep_vote=3]
shared is_ready: prepared -> ready
shared start_voting: ready -> voting
shared aID_vote: voting -> vote [Coercer1_aID_vote=?aID_vote, Coercer1_aID_revote=?aID_revote]
shared send_vote_aID: vote -> send
revote_vote_1: send -[aID_revote==1]> voting [aID_vote=?aID_required, aID_revote=2]
skip_revote_1: send -[aID_revote==1]> votingf
revote_vote_2: send -[aID_revote==2]> voting [aID_vote=?aID_required, aID_revote=3]
skip_revote_2: send -[aID_revote==2]> votingf
final_vote: send -[aID_revote==3]> votingf [aID_vote=?aID_prep_vote]
skip_final: send -[aID_revote==3]> votingf
shared send_fvote_aID: votingf -> sendf
shared finish_voting: sendf -> finish
shared send_tracker_aID: finish -> tracker
shared finish_sending_trackers: tracker -> trackers_sent
shared give1_aID: trackers_sent -> interact [Coercer1_aID_tracker=1]
shared give2_aID: trackers_sent -> interact [Coercer1_aID_tracker=2]
shared not_give_aID: trackers_sent -> interact [Coercer1_aID_tracker=0]
shared punish_aID: interact -> ckeck [aID_punish=true]
shared not_punish_aID: interact -> check [aID_punish=false]
shared check_tracker1_aID: check -> end
shared check_tracker2_aID: check -> end
PROTOCOL: [[coerce1_aID, coerce2_aID, coerce3_aID], [punish, not_punish]]
```

# Formula

$$\varphi_{vuln,i,k} = \langle\langle Coercer \rangle\rangle G((end \wedge revote_{v_1} = k \wedge voted_{v_1} = i) \rightarrow K_{Coercer} voted_{vi} = i)$$

## Configurations:

- First candidate ( $i = 1$ ) and  $k = \#R$  revotes
- Last candidate ( $i = \#C$ ) and  $k = \#R$  revotes
- First candidate ( $i = 1$ ) and  $k = \#R - 1$  revotes
- Last candidate ( $i = \#C$ ) and  $k = \#R - 1$  revotes

#A	#R	Full Model					Reduced Model					Result
		#st	#tr	Seq.	Par.	Appr.	#st	#tr	Seq.	Par.	Appr.	
4	3	3.63e4	7.46e4	0.003	0.009	1.121	2.60e4	5.99e4	0.001	0.002	0.184	True
4	5	5.62e4	1.15e5	0.004	0.003	0.345	4.01e4	9.26e4	0.002	0.002	0.283	True
4	10	1.06e5	2.18e5	0.009	0.005	0.691	7.55e4	1.74e5	0.004	0.002	0.563	True
5	3	1.55e6	5.91e6	0.158	0.004	14.78	1.09e6	4.65e6	0.112	0.021	12.99	True
6	3	7.61e7	4.98e8	0.524	0.051	41.24	5.34e7	3.82e8	0.427	0.042	37.35	True
7	3	model generation timeout					model generation timeout					-

Table 1

Verification of  $\varphi_{vuln,i,k}$  for the first candidate ( $i = 1$ ) and  $k = \#R$  revotes

#Ag	#R	Full Model					Reduced Model					Result
		#st	#tr	Seq.	Par.	Appr.	#st	#tr	Seq.	Par.	Appr.	
4	3	3.63e4	7.46e4	0.003	0.010	1.103	2.60e4	5.99e4	0.002	0.003	0.166	True
4	5	5.62e4	1.15e5	0.004	0.005	0.348	4.01e4	9.26e4	0.003	0.003	0.280	True
4	10	1.06e5	2.18e5	0.008	0.009	0.700	7.55e4	1.74e5	0.005	0.004	0.567	True
5	3	1.55e6	5.91e6	0.160	0.055	14.03	1.09e6	4.65e6	0.112	0.053	12.49	True
6	3	7.61e7	4.98e8	0.602	0.083	42.44	5.34e7	3.82e8	0.501	0.057	38.20	True
7	3	model generation timeout					model generation timeout					-

Table 2

Verification of  $\varphi_{vuln,i,k}$  for the last candidate ( $i = \#C$ ) and  $k = \#R$  revotes

#Ag	#R	Full Model					Reduced Model					Result
		#st	#tr	Seq.	Par.	Appr.	#st	#tr	Seq.	Par.	Appr.	
4	3	3.63e4	7.46e4	0.303	0.317	1.128	2.60e4	5.99e4	0.202	0.205	0.179	False
4	5	5.62e4	1.15e5	0.524	0.592	0.325	4.01e4	9.26e4	0.411	0.503	0.280	False
4	10	1.06e5	2.18e5	0.721	0.668	0.459	7.55e4	1.74e5	0.525	0.512	0.364	False
5	3	1.55e6	5.91e6	2.146	1.257	0.981	1.09e6	4.65e6	1.513	1.003	0.583	False
6	3	7.61e7	4.98e8	5.232	3.228	1.892	5.34e7	3.82e8	4.986	2.427	1.092	False
7	3	model generation timeout					model generation timeout					-

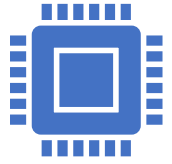
Table 3

Verification of  $\varphi_{vuln,i,k}$  for the first candidate ( $i = 1$ ) and  $k = \#R - 1$  revotes

#Ag	#R	Full Model					Reduced Model					Result
		#st	#tr	Seq.	Par.	Appr.	#st	#tr	Seq.	Par.	Appr.	
4	3	3.63e4	7.46e4	0.302	0.311	0.180	2.60e4	5.99e4	0.201	0.213	0.126	False
4	5	5.62e4	1.15e5	0.519	0.584	0.310	4.01e4	9.26e4	0.410	0.475	0.283	False
4	10	1.06e5	2.18e5	0.742	0.627	0.462	7.55e4	1.74e5	0.558	0.544	0.370	False
5	3	1.55e6	5.91e6	2.160	1.358	0.942	1.09e6	4.65e6	1.621	1.009	0.519	False
6	3	7.61e7	4.98e8	5.504	3.516	1.903	5.34e7	3.82e8	5.110	2.380	1.112	False
7	3	model generation timeout					model generation timeout					-

Table 4

Verification of  $\varphi_{vuln,i,k}$  for the last candidate ( $i = \#C$ ) and  $k = \#R - 1$  revotes



# Results

- DominoDFS and alternative distributed algorithm performed much slower and are omitted from the results
- Parallel verification performs quite well in most cases
- Performance of the parallel algorithm depends heavily on the structure of the model
- The fixpoint approximation performs well in cases where no strategy can be found



# Conclusions



Modal logics for MAS are characterized by high computational complexity.



We used the „all out” approach, verifying a genuine protocol for secure voting.



Partial-order reductions, simple DFS, simple distributed DFS and fixpoint approximation show very promising performance.



Thank you for  
your  
attention!